Tradeoff to minimize extra-computations and stopping criterion tests for parallel iterative schemes

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In general, iterative schemes consist in iterating some computation until a certain boolean condition is carried out. Unfortunately, for the parallel implementation on distributed memory systems, the stopping criterion test is often achieved by global synchronization which penalizes performances, in particular in the context of grid computing.

In order to mitigate this problem, different approaches have been proposed and studied in the literature. A first approach consists in gathering all global synchronizations of one iteration in order to perform them together; this technique is used in [1] for Conjugate Gradient method, but it does not allow to reduce the number of stopping criterion tests. A second approach, which has been widely studied, consists in desynchronizing communications and iterations, consequently resulting in an asynchronous algorithm[2, 3]. Unfortunately, in this context, the iteration scheme is no more preserved; thus, a critical problem is to ensure convergence. A third approach, called \(k\)-steps in the sequel, is often used in practice; it consists in performing the stopping tests every \(k\) iterations only (i.e. stopping tests are evaluated at iterations \(k, 2k, \ldots, qk, \ldots\)), till the stopping test is carried out. Let \(n_s\) denotes the exact number of iterations obtained by the standard algorithm which performs the stopping criterion test after each iteration. Then, the \(k\)-steps method performs \(#I = n_s + k\) iterations and \(#T = [n_s/k]\) tests. However, since \(n_s\) is unknown a priori, the problem is to choose \(k\) in order to obtain a good trade-off between \(#I\) and \(#T\).

In this work, we propose a new control technique, original to the best of our knowledge, which delivers the result of iteration \(n_s\) after a small number \(#T = \log^{1+o(1)} n_s\) of stopping criterion tests while computing the result of \(#I = n_s(1 + e)\) iterations only. The advantage of our technique lies in the fact that the structure of the initial algorithm is not modified, and consequently the convergence is ensured.

Our method is based on an amortized technique inspired from Floyd’s and Brent’s algorithms to detect periodicity in a sequence [4]. It consists in computing two numbers \(n_1\) and \(n_2\) such that \(n_1 \leq n_s \leq n_2\), and then determining the exact value of \(n_s\). Like the \(k\)-steps, our method performs only a limited number of stopping tests, however tests are not performed at regular steps anymore, but rather at steps \(\rho^{f(0)}, \rho^{f(1)}, \ldots, \rho^{f(k)} = n_2\), where \(1 < \rho < 2\) and \(f\) satisfies \(\forall i, f(i) \leq i\). Determination of \(f\) is based on a trade-off between \(#I\) and \(#T\). In this paper, we provide asymptotic results for different choices of \(f\): namely \(f(i) = i^\alpha, 0 < \alpha \leq 1\); \(f(i) = \log i\), and \(f(i) = \frac{1}{\log i}\). We prove that \(n_s\) can be determined with very little more than \(#T = \log n_s\) tests – which is a lower bound for \(#T\) – while performing an asymptotic optimal number of iterations \(#I = n_s + o(n_s)\).

References


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